

Introduction to E-Graphs

Rebecca Swords | Women in Compilers and Tools

Questions

What are e-graphs?

What are they good for?

How do they work?

E-Graph

*A data structure representing an
equivalence relation over terms*

Practical Applications

Theorem proving

SMT solving

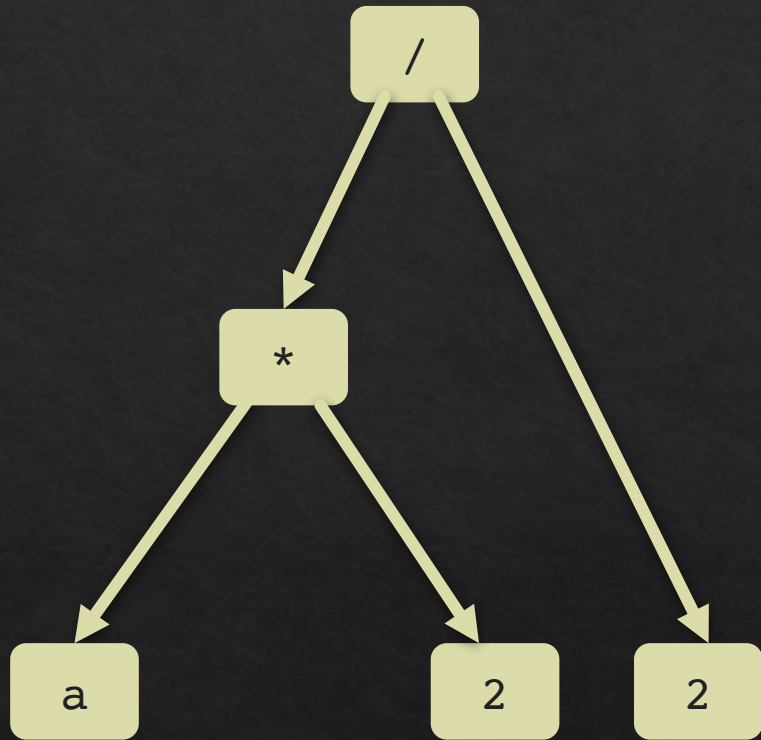
Optimization

Translation validation

Compilation

Synthesis

Running Example: $(a * 2) / 2$

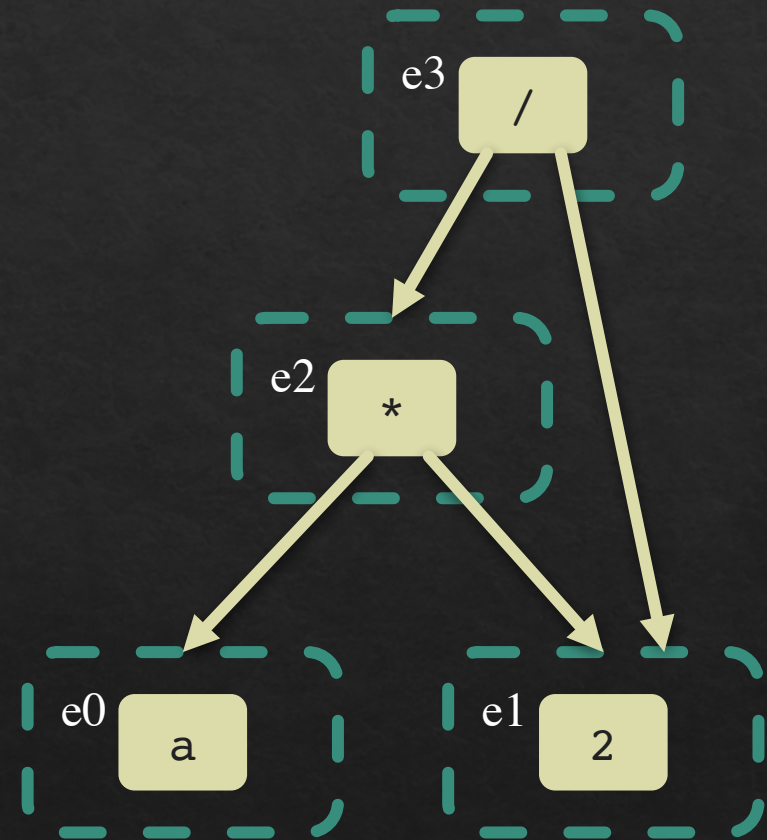
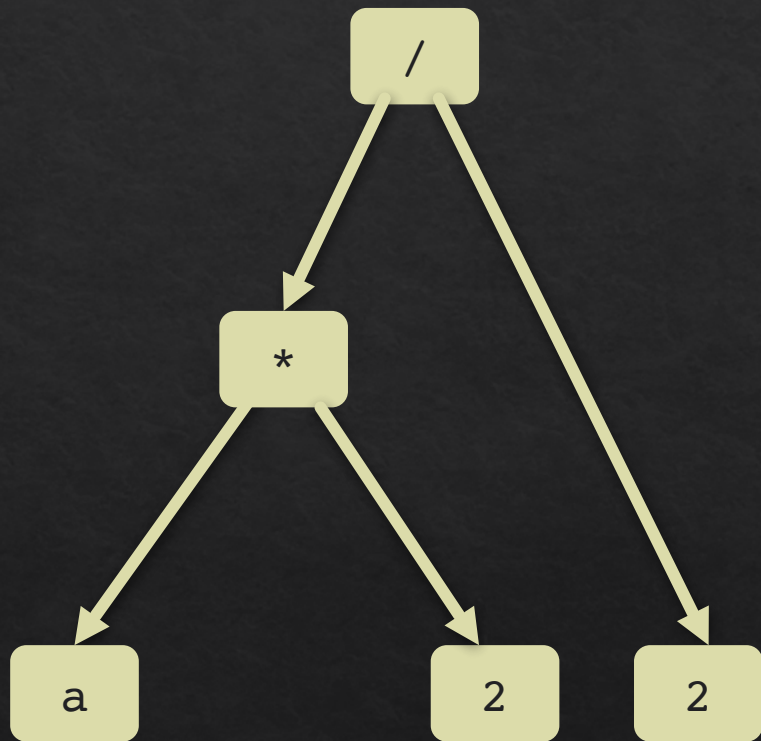


→ This reduces to a

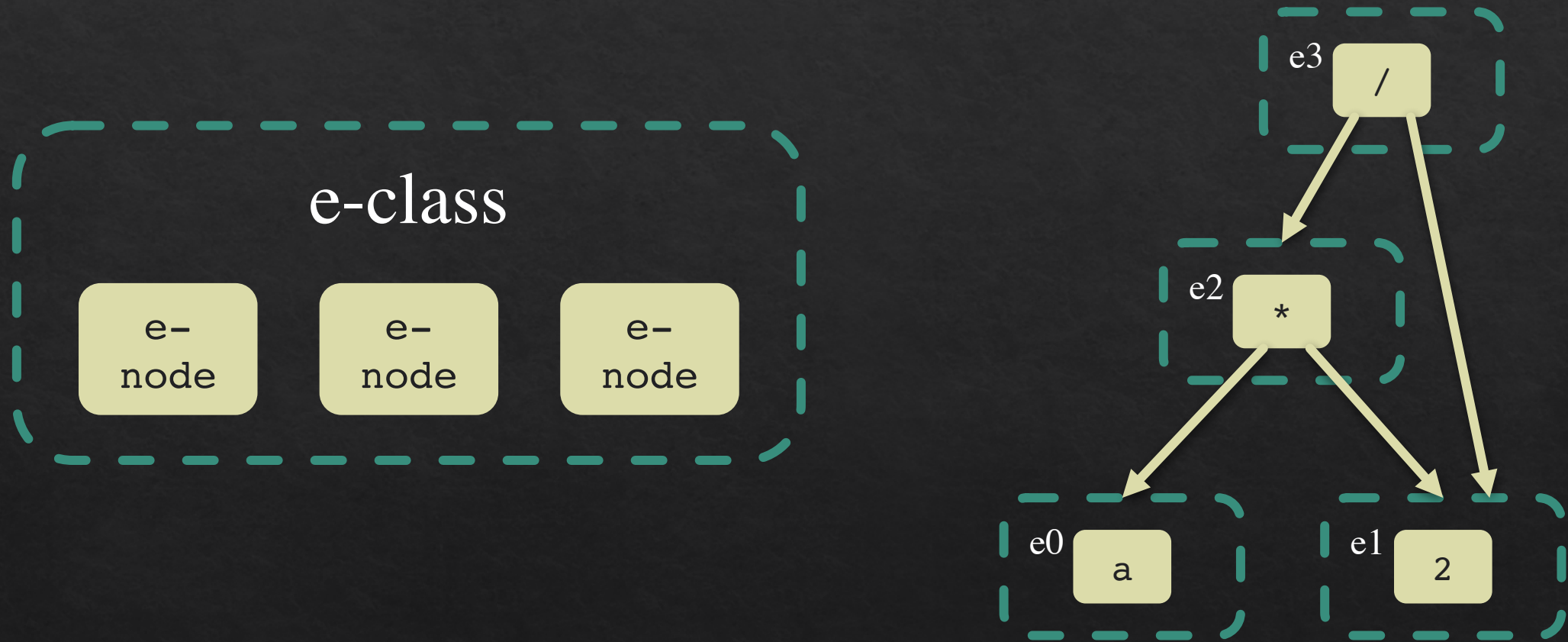
→ We can use e-graphs to do it!

Start with this AST

Into an E-Graph



Into an E-Graph

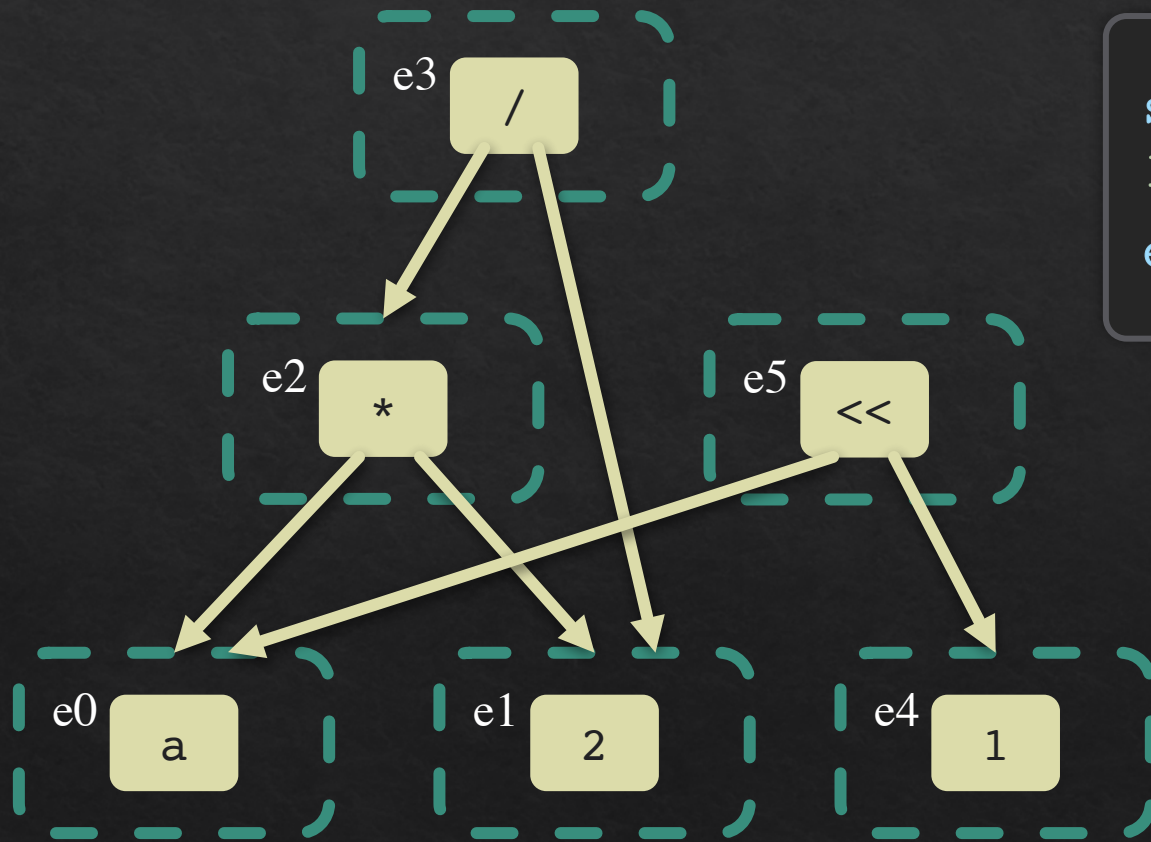


Build it with Quiche

```
1 expr = (ExprNode('a', ())) * 2) /  
2  
2 quiche_tree = ExprTree(expr)  
3 egraph = EGraph(quiche_tree)
```

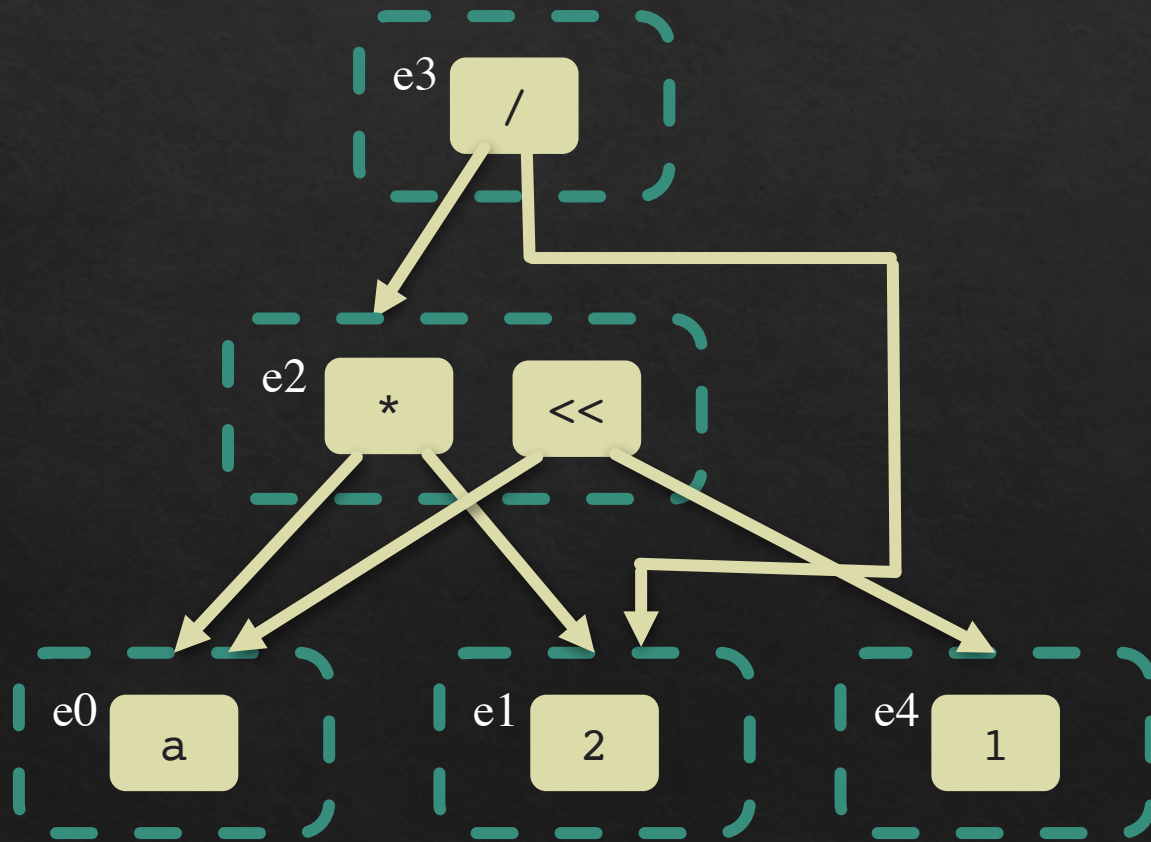
1. Parse term into the arithmetic language structure
2. Construct intermediate QuicheTree representation
3. Create e-graph from QuicheTree

Another Term: $a \ll 1$



```
shift_expr = ExprNode('a', ()) <<  
1  
egraph.add(ExprTree(shift_expr))
```

Merging Equivalent Terms



We assert:

$$a * 2 === a \ll 1$$

It follows that:

$$(a * 2) / 2 === (a \ll 1) / 2$$

Manual Merging in Quiche

```
1 shift_eclass =  
  egraph.add(ExprTree(shift_expr))  
  times_node = ExprNode('a', ()) * 2  
  times_eclass =  
    egraph.add(ExprTree(times_node))  
2 egraph.merge(times_eclass, shift_eclass)  
3 egraph.rebuild()
```

1. Save e-class IDs for the expressions to be merged
2. Merge the two e-classes together
3. Restore e-graph invariants

E-Graphs More Formally

Structure

- ◇ E-node: an n-ary function symbol and n children (e-class IDs)
- ◇ E-class: set of e-nodes
- ◇ Union-find over e-classes: `add`, `merge`, `find` operations
- ◇ Canonical e-node: for each child, `i`, `find(i) = i`
- ◇ Hashcons: maps canonical e-nodes to e-classes

Invariants

- ◇ Hashcons maps all canonical e-nodes
- ◇ Equivalence closed under congruence, i.e., congruent e-nodes are in the same e-class
If $a = b$, then $f(a) = f(b)$

Why is this good for term rewriting?

Instead of destructive rewrites, put *all* equivalent terms in the e-graph

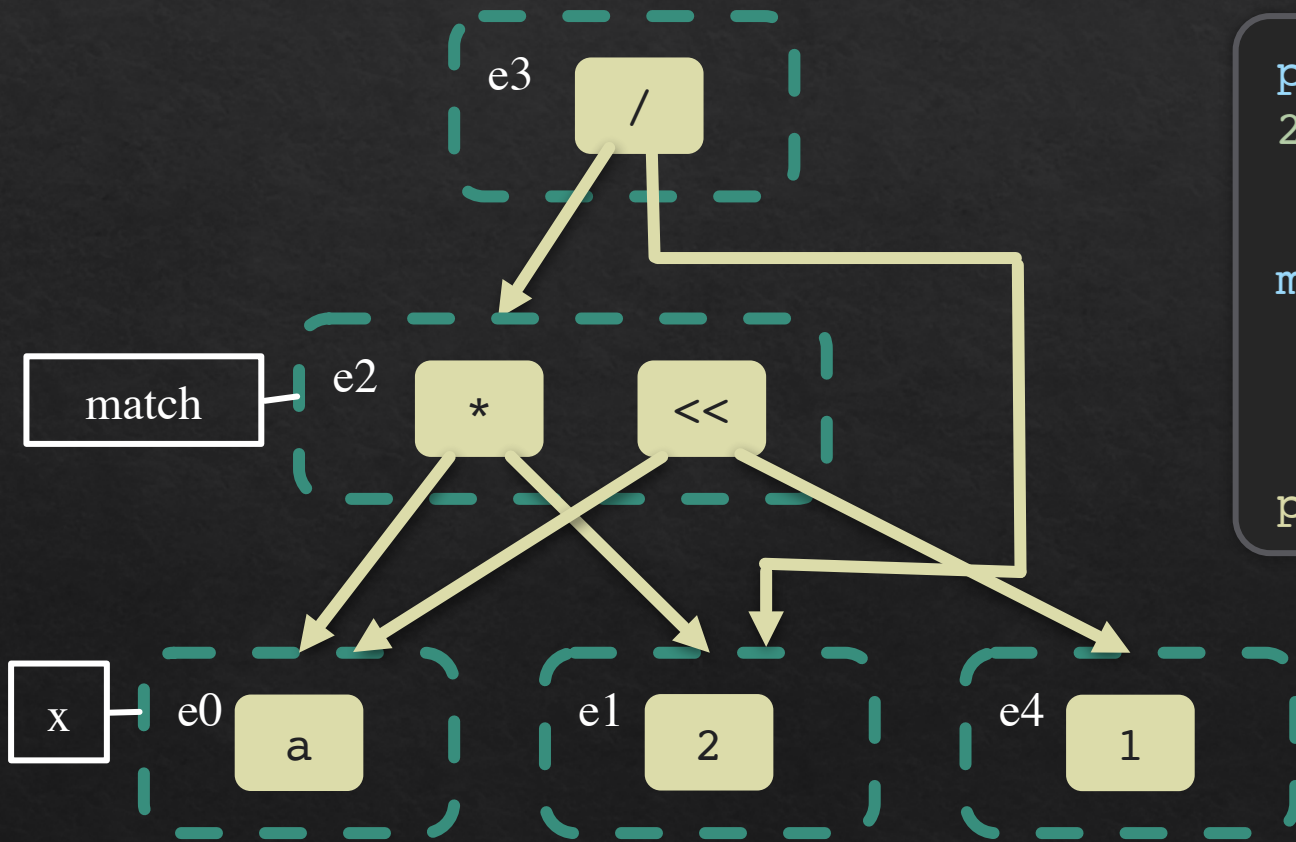
- No worries about phase ordering
- Consider all options and choose the “best” at the end

E-Matching

*Pattern matching
for e-graphs!*

- Add pattern variables to language
- `ematch` searches for a pattern and returns:
 - ◇ e-class matching the term
 - ◇ substitution from vars to e-class IDs

E-Matching Example: $x * 2$



```
pattern = ExprTree(ExprNode('x', ()) *  
2)
```

```
matches = egraph.ematch(pattern,  
    egraph.eclases())
```

```
print(matches)
```

```
[(e2, { 'x': e0 })]
```

Rewriting Rules: Pattern Merges

1

```
rule = ExprTree.make_rule(lambda x:  
    (x * 2, x << 1))
```

2

```
Rule.apply_rules([rule], egraph)
```

3

```
print("Shift e-class: ", shift_eclass)  
print("Shift e-class.find(): ",  
      shift_eclass.find())
```

1. Create a rule:

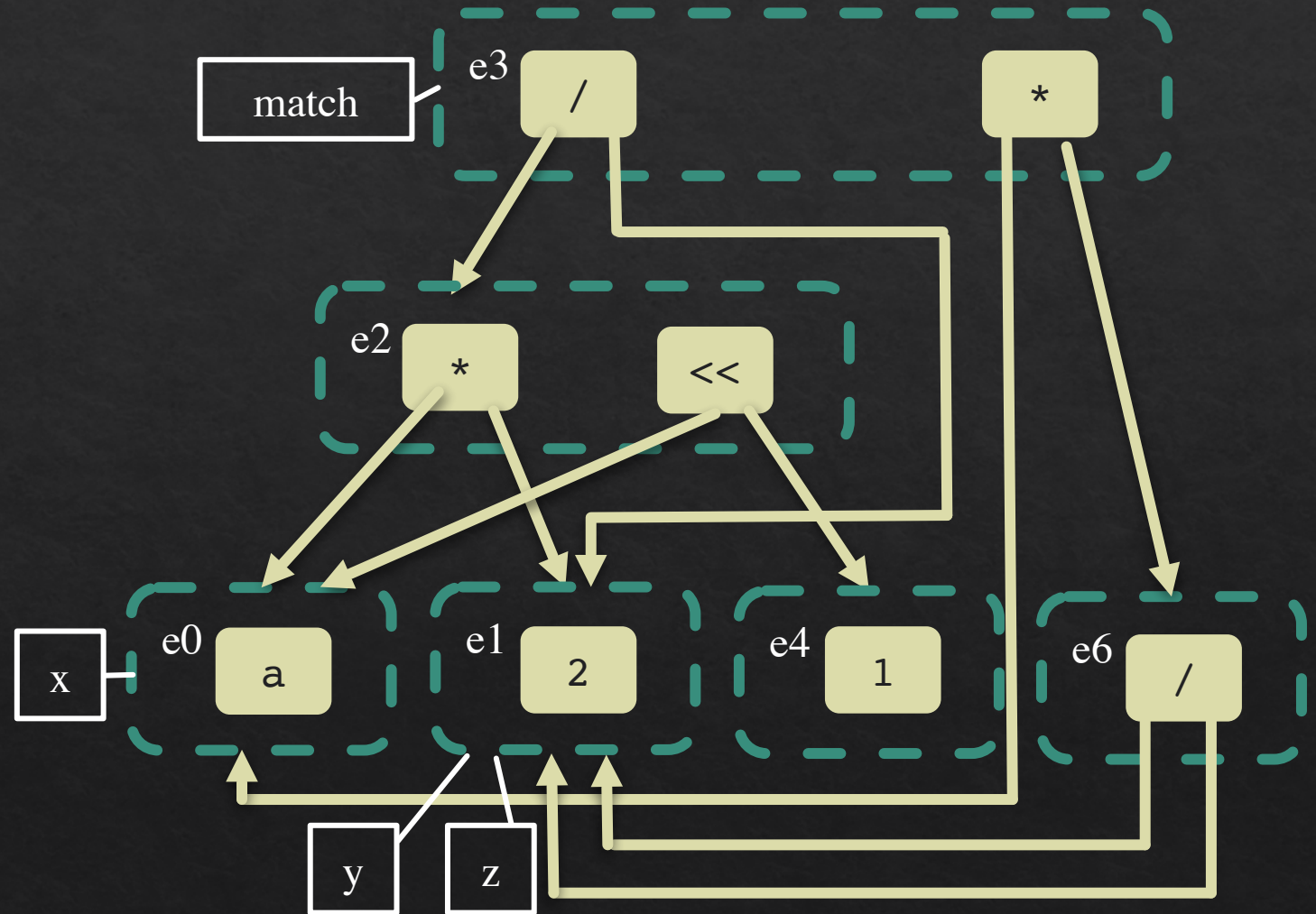
$x * 2 == x \ll 1$

2. Apply all rules to e-graph (and rebuild)

3. Shift e-class: e5
Shift e-class find: e2

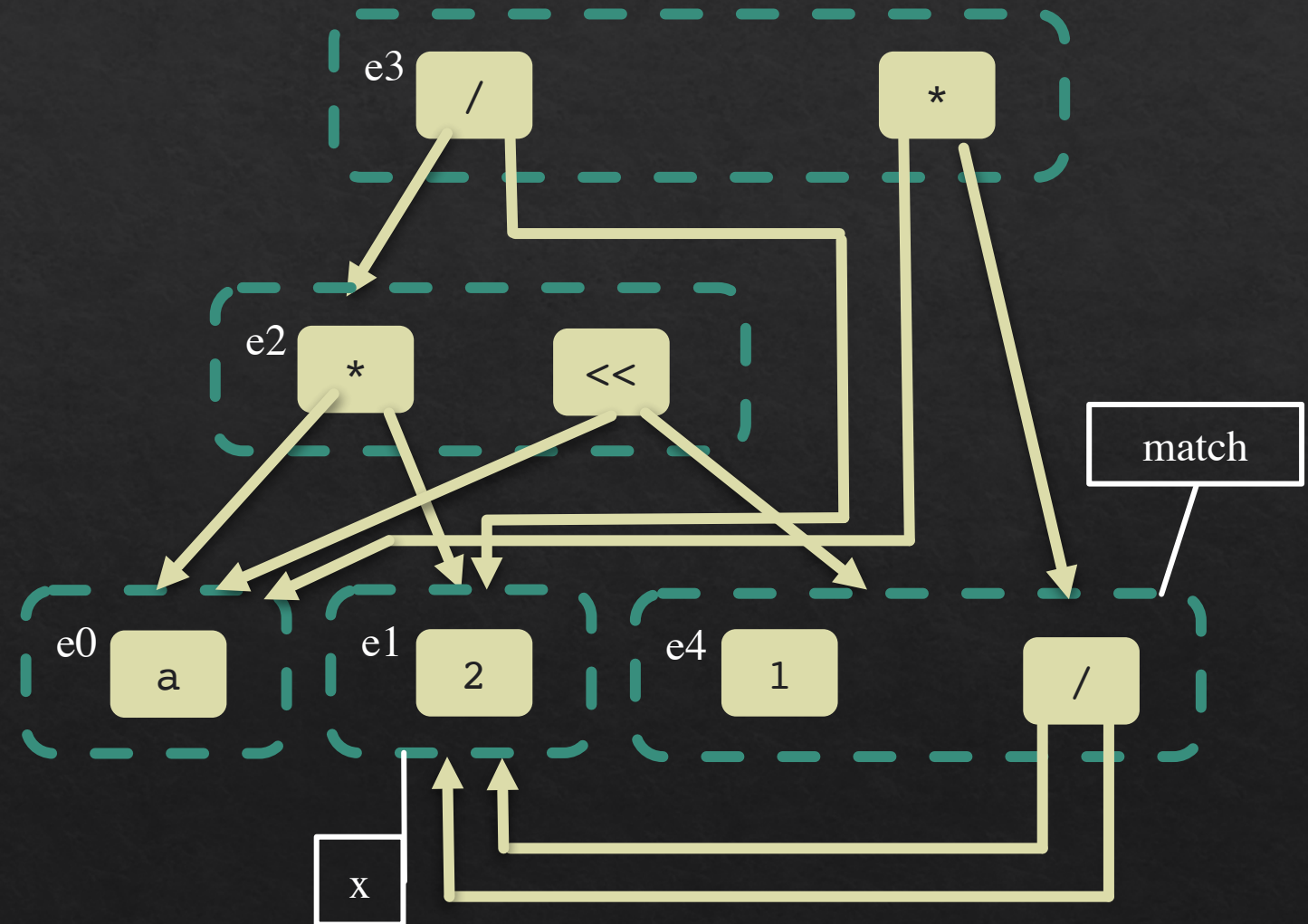
Another rewrite:

$$\begin{aligned} (x * y) / z \\ === \\ x * (y / z) \end{aligned}$$



And another:

x/x
====
1



And one more:

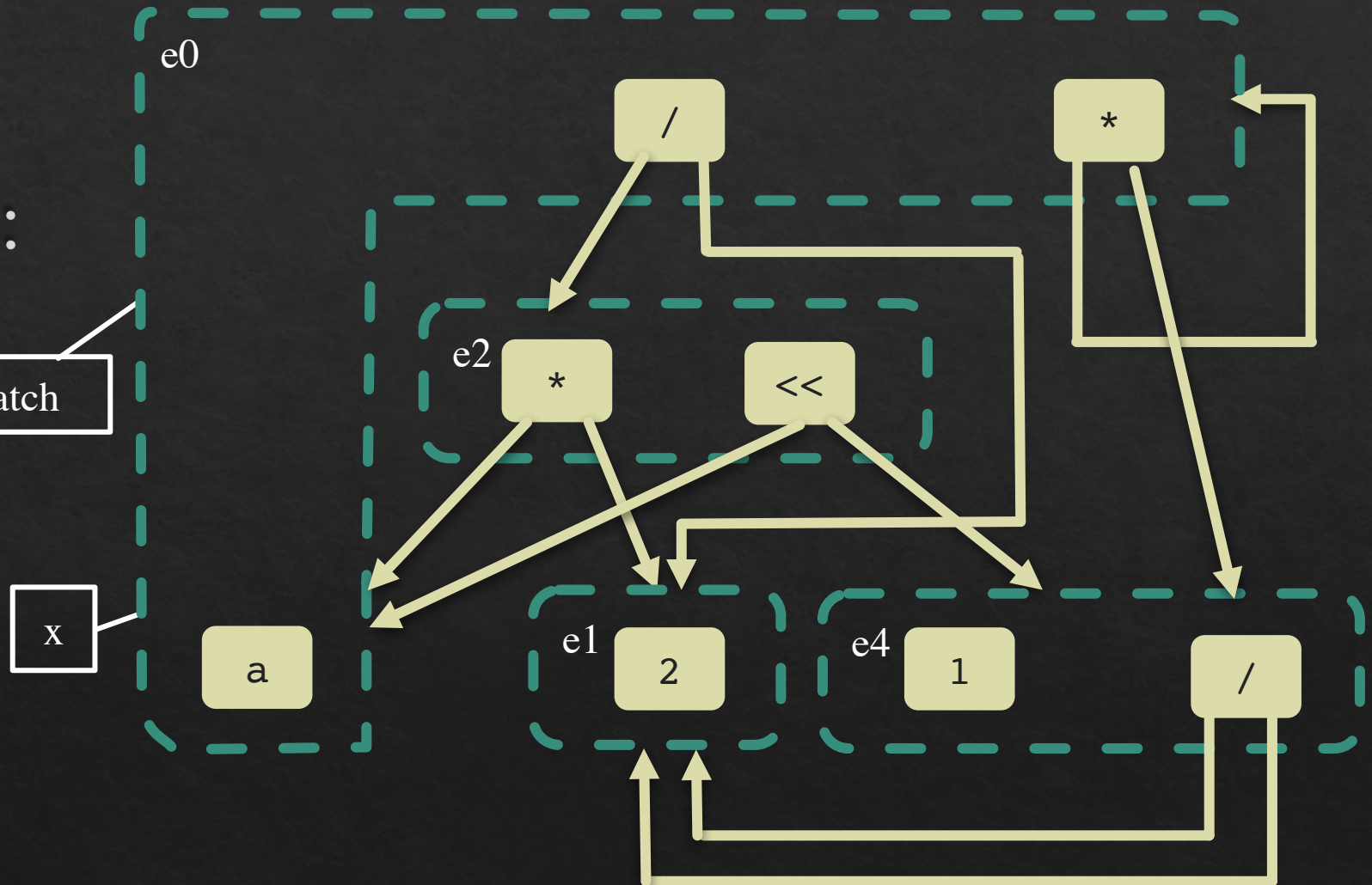
$x * 1$

====

x

match

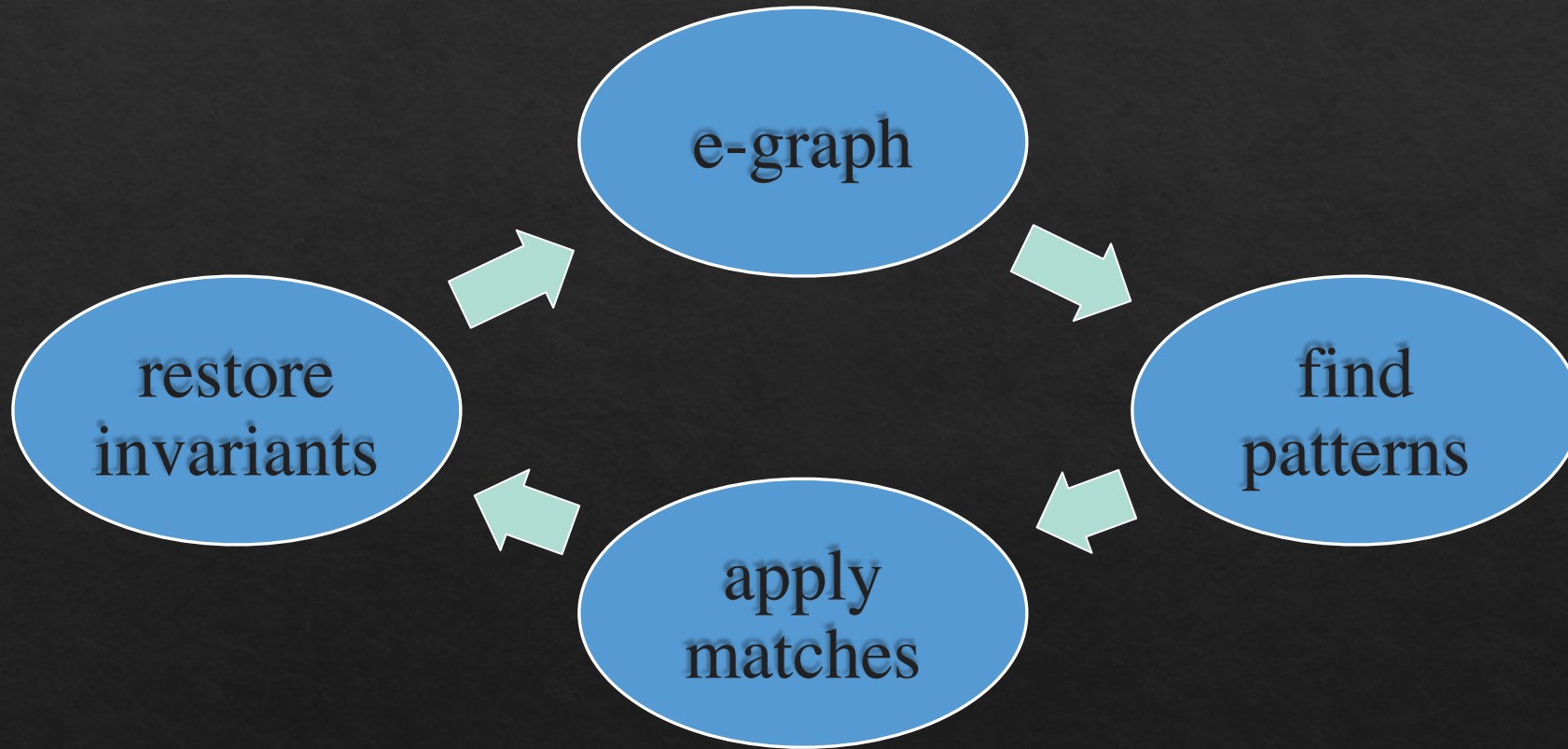
x



Keep applying rewrite
rules until no new
changes are made

Equality
Saturation

Equality Saturation Loop



Apply Rules Until Saturation

1

```
rules = [  
    ExprTree.make_rule(lambda x, y, z:  
        ((x * y) / z, x * (y / z))),  
    ExprTree.make_rule(lambda x:  
        (x / x, ExprNode(1, ()))),  
    ExprTree.make_rule(lambda x: (x * 1, x))  
]
```

2

```
while not egraph.is_saturated():  
    Rule.apply_rules(rules, egraph)  
aeclass = egraph.add(ExprTree(ExprNode('a',  
    ())))  
assert aeclass.find() == egraph.root.find()
```

3

1. Same 3 rules we just applied
2. Apply rules until the e-graph is saturated
3. Verification: expect a to have merged with the “root” e-class

E-Class Analysis

*Domain-specific e-
graph extensions*

- ◆ Attach datum to each e-class based on e-nodes: `make`
- ◆ Merge data when e-classes merge: `join`
- ◆ Update e-class based on datum: `modify`
- ◆ Form a join-semilattice

What Can E-Class Analyses Do?

- ◇ Program analysis
- ◇ Conditional or dynamic rewrites
- ◇ Debugging
- ◇ Pruning
- ◇ On-the-fly term extraction

Standardized interface for extending e-graphs!

Analysis Invariant

for each e-class

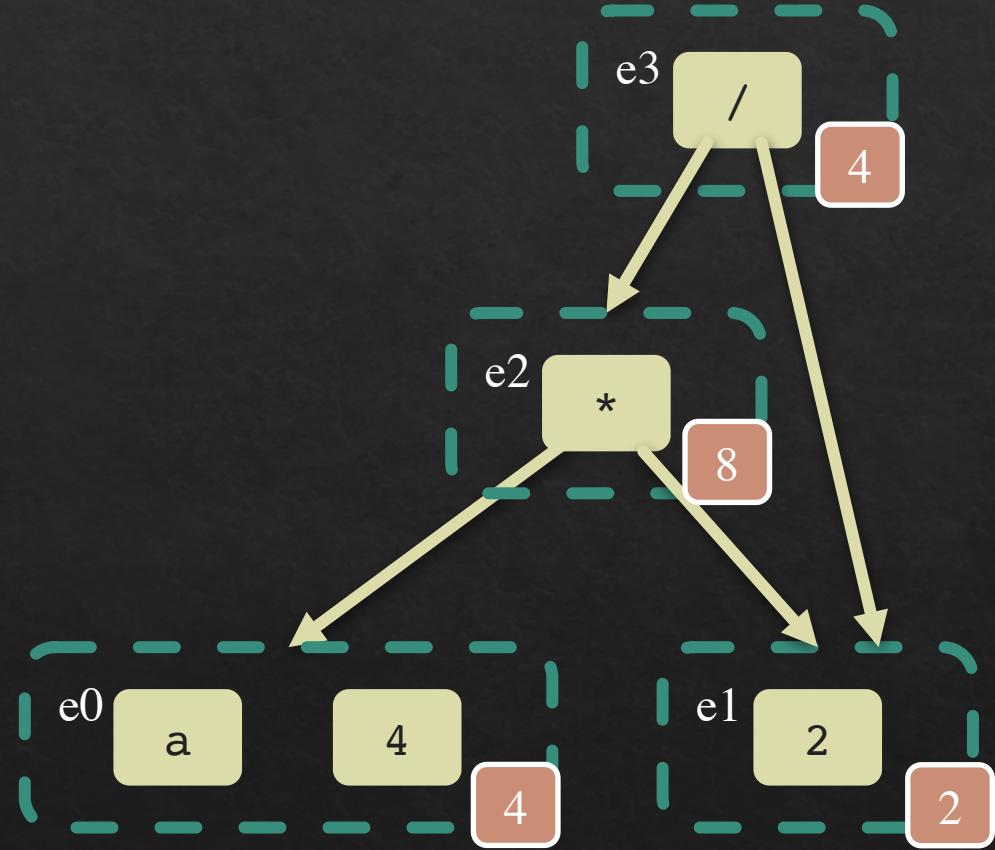
fixed point

$$\forall c \in G. \quad d_c = \bigwedge_{n \in c} \text{make}(n) \quad \text{and} \quad \text{modify}(c) = c$$

data is the same as `make`-ing
data for each e-node and then
`join`-ing

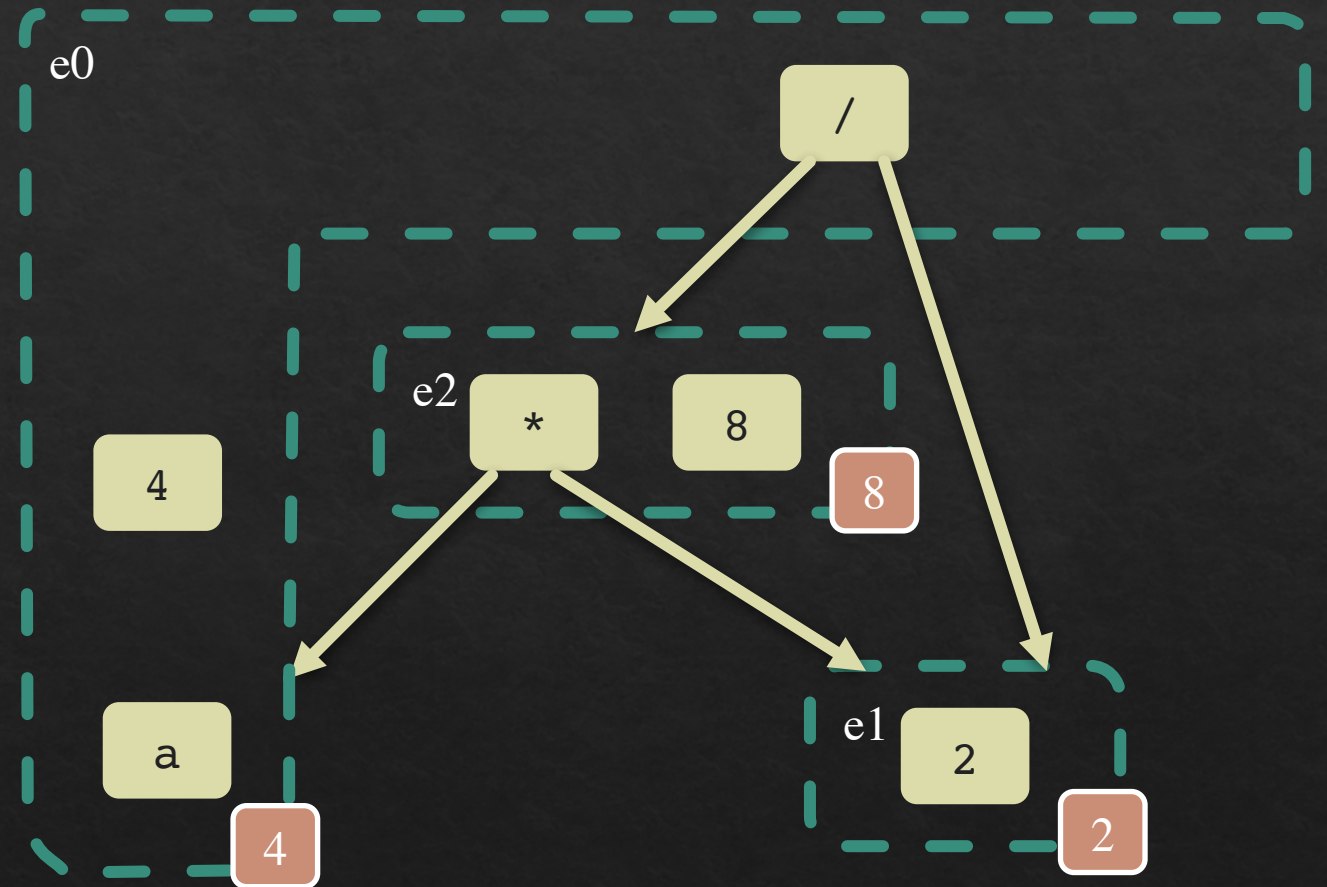
Constant Folding E-Class Analysis

Suppose we learn that a
=== 4



Constant Folding E-Class Analysis

Suppose we learn that a
=== 4



Constant Folding: Usage

1

```
expr = (ExprNode('a', ()) * 2) / 2
```

```
quiche_tree = ExprTree(expr)
```

```
egraph = EGraph(quiche_tree, ExprConstantFolding())
```

2

```
four_eclass = egraph.add(ExprTree(ExprNode(4, ())))
```

```
a_eclass = egraph.add(ExprTree(ExprNode("a", ())))
```

3

```
egraph.merge(a_eclass, four_eclass)
```

4

```
egraph.rebuild()
```

5

```
assert egraph.root.data == 4
```

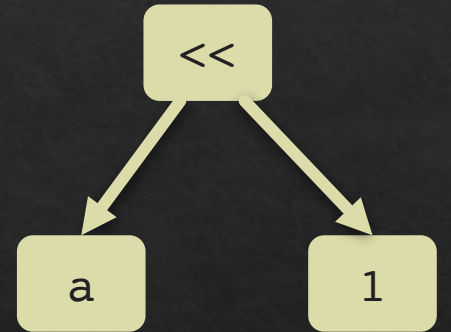
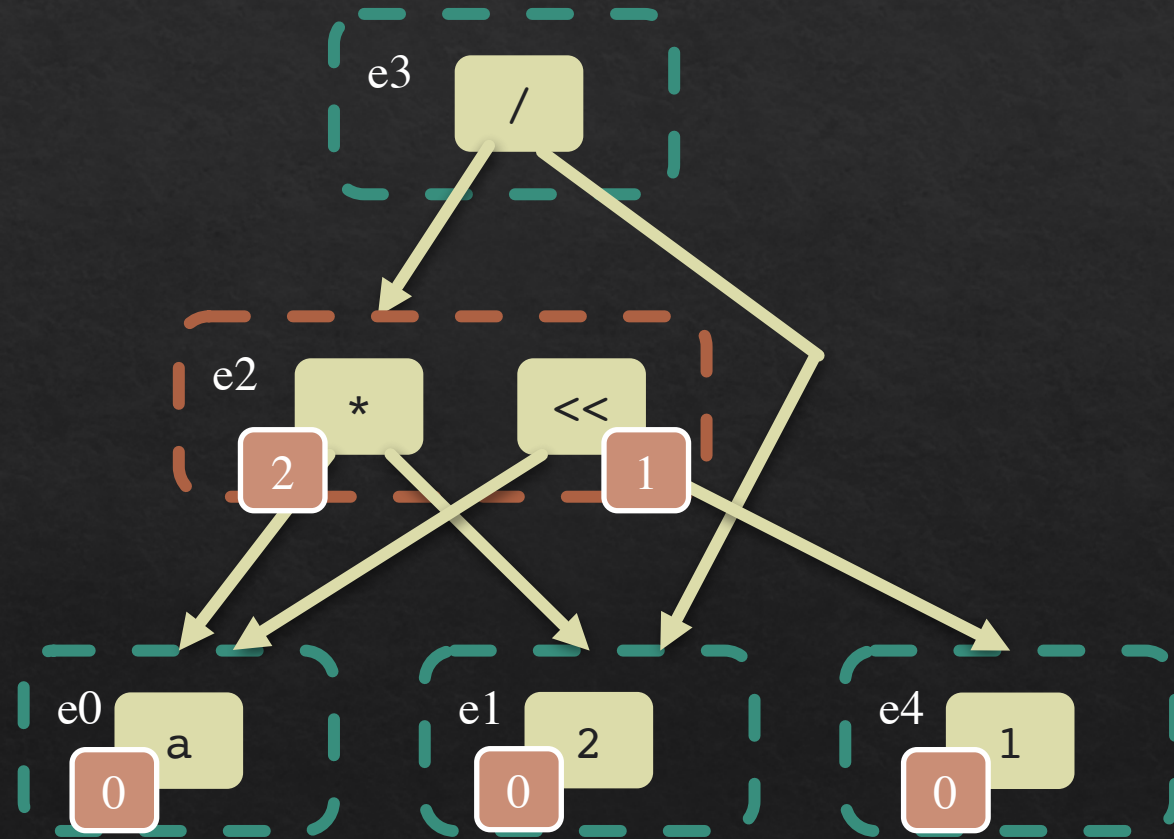
1. Create e-graph with constant folding analysis
2. Get e-class IDs
3. Merge 4 with a
4. Rebuild (update analysis)
5. Verify

Term Extraction

- ◆ Pick an e-class to extract
- ◆ Cost model assigns a cost to e-nodes
- ◆ Choose best e-node for each e-class
- ◆ Construct a term by combining the e-node values

Term Extraction Example

Operator	Cost
+	1
<<	1
*	2
/	3
default	0



Term Extraction Example

```
1 cost_model = ExprNodeCost()  
2 extractor= MinimumCostExtractor()  
3 extracted = extractor.extract(  
4     cost_model,  
    egraph,  
    egraph.root.find(),  
    ExprTree.make_node)  
assert str(extracted) == "a"
```

1. Initialize cost model and extractor
2. Extract the best term
3. Specify which e-class to extract
4. Function to construct ExprTree from e-node data

More on Quiche

- ◆ Add your own languages!
 - ◆ Bring your own parser, adapt your AST into a QuicheTree
- ◆ End-to-end Python rewriting!
 - ◆ Uses native Python parser (v3.7+)
 - ◆ Read/write valid Python files
- ◆ Native Python!
 - ◆ With all its pros and cons

QuicheTree

Quiche requires the user to provide a parsed tree that implements `QuicheTree` (“*bring your own parser*”).

value()

the e-node key

children()

list of the node's children

is_pattern_symbol()

for e-matching; indicates if the node is a pattern

```
class QuicheTree(ABC):  
    @abstractmethod  
    def value(self)  
  
    @abstractmethod  
    def children(self)  
  
    @abstractmethod  
    def is_pattern_symbol(self)
```

Links and References

- ◇ Quiche repo: <https://github.com/riswords/quiche>
- ◇ egg website: <https://egraphs-good.github.io/>
- ◇ egg: Fast and extensible equality saturation (POPL '21, Willsey, et al.): <https://dl.acm.org/doi/10.1145/3434304>
- ◇ Equality-Based Translation Validator for LLVM (CAV '11, Stepp, Tate, & Lerner): https://cseweb.ucsd.edu/~rtate/publications/eqsat/eqsat_stepp_cav11.pdf
- ◇ babble: Learning Better Abstractions with E-Graphs and Anti-Unification (POPL '23, Cao, et al.): <https://dl.acm.org/doi/10.1145/3571207>



Questions?

Additional References from Q&A

1. Link to the public E-Graphs Zulip chat: <https://egraphs.zulipchat.com/>
2. Perfect Reconstructability of Control Flow from Demand Dependence Graphs (Bahmann, et al. 2014) <https://dl.acm.org/doi/abs/10.1145/2693261>
3. E-Graphs Zulip discussion of using RVSDG representation: <https://egraphs.zulipchat.com/#narrow/stream/328976-Program-Optimization/topic/PEGs>
4. Equality Saturation for Tensor Graph Superoptimization (Yang, et al., MLSys 2014): <https://arxiv.org/abs/2101.01332>
5. Relational e-matching (Zhang, et al., POPL 2022)
6. [Logging an Egg: Datalog on E-Graphs \(EGRAPHS 2022\) - PLDI 2022 \(sigplan.org\)](#)